

MCP2510

Stand-Alone CAN Controller with SPI™ Interface

Features

- Implements Full CAN V2.0A and V2.0B at 1 Mb/s:
	- 0 8 byte message length
	- Standard and extended data frames
	- Programmable bit rate up to 1 Mb/s
	- Support for remote frames
	- Two receive buffers with prioritized message storage
	- Six full acceptance filters
	- Two full acceptance filter masks
	- Three transmit buffers with prioritization and abort features
	- Loop-back mode for self test operation
- Hardware Features:
	- High Speed SPI Interface (5 MHz at 4.5V I temp)
	- Supports SPI modes 0,0 and 1,1
	- Clock out pin with programmable prescaler
	- Interrupt output pin with selectable enables
	- 'Buffer full' output pins configureable as interrupt pins for each receive buffer or as general purpose digital outputs
	- 'Request to Send' input pins configureable as control pins to request immediate message transmission for each transmit buffer or as general purpose digital inputs
	- Low Power Sleep mode
- Low power CMOS technology:
	- Operates from 3.0V to 5.5V
	- 5 mA active current typical
	- 10 µA standby current typical at 5.5V
- 18-pin PDIP/SOIC and 20-pin TSSOP packages
- Temperature ranges supported:
	- Industrial (I): -40° C to $+85^{\circ}$ C

Description

The Microchip Technology Inc. MCP2510 is a Full Controller Area Network (CAN) protocol controller implementing CAN specification V2.0 A/B. It supports CAN 1.2, CAN 2.0A, CAN 2.0B Passive, and CAN 2.0B Active versions of the protocol, and is capable of transmitting and receiving standard and extended messages. It is also capable of both acceptance filtering and message management. It includes three transmit buffers and two receive buffers that reduce the amount of microcontroller (MCU) management required. The MCU communication is implemented via an industry standard Serial Peripheral Interface (SPI) with data rates up to 5 Mb/s.

Package Types

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1.0 DEVICE FUNCTIONALITY

1.1 Overview

The MCP2510 is a stand-alone CAN controller developed to simplify applications that require interfacing with a CAN bus. A simple block diagram of the MCP2510 is shown in [Figure 1-1](#page-2-1). The device consists of three main blocks:

- 1. The CAN protocol engine.
- 2. The control logic and SRAM registers that are used to configure the device and its operation.
- 3. The SPI protocol block.

A typical system implementation using the device is shown in [Figure 1-2.](#page-3-0)

The CAN protocol engine handles all functions for receiving and transmitting messages on the bus. Messages are transmitted by first loading the appropriate message buffer and control registers. Transmission is initiated by using control register bits, via the SPI interface, or by using the transmit enable pins. Status and errors can be checked by reading the appropriate registers. Any message detected on the CAN bus is

FIGURE 1-1: BLOCK DIAGRAM

checked for errors and then matched against the user defined filters to see if it should be moved into one of the two receive buffers.

The MCU interfaces to the device via the SPI interface. Writing to and reading from all registers is done using standard SPI read and write commands.

Interrupt pins are provided to allow greater system flexibility. There is one multi-purpose interrupt pin as well as specific interrupt pins for each of the receive registers that can be used to indicate when a valid message has been received and loaded into one of the receive buffers. Use of the specific interrupt pins is optional, and the general purpose interrupt pin as well as status registers (accessed via the SPI interface) can also be used to determine when a valid message has been received.

There are also three pins available to initiate immediate transmission of a message that has been loaded into one of the three transmit registers. Use of these pins is optional and initiating message transmission can also be done by utilizing control registers accessed via the SPI interface.

[Table 1-1](#page-3-1) gives a complete list of all of the pins on the MCP2510.

TABLE 1-1: PIN DESCRIPTIONS

Note: Type Identification: I=Input; O=Output; P=Power

1.2 Transmit/Receive Buffers

The MCP2510 has three transmit and two receive buffers, two acceptance masks (one for each receive buffer), and a total of six acceptance filters. [Figure 1-3](#page-4-0) is a block diagram of these buffers and their connection to the protocol engine.

1.3 CAN Protocol Engine

The CAN protocol engine combines several functional blocks, shown in [Figure 1-4](#page-5-0). These blocks and their functions are described below.

1.4 Protocol Finite State Machine

The heart of the engine is the Finite State Machine (FSM). This state machine sequences through messages on a bit by bit basis, changing states as the fields of the various frame types are transmitted or received. The FSM is a sequencer controlling the sequential data stream between the TX/RX Shift Register, the CRC Register, and the bus line. The FSM also controls the Error Management Logic (EML) and the parallel data stream between the TX/RX Shift Registers and the buffers. The FSM insures that the processes of reception, arbitration, transmission, and error signaling are performed according to the CAN protocol. The automatic retransmission of messages on the bus line is also handled by the FSM.

1.5 Cyclic Redundancy Check

The Cyclic Redundancy Check Register generates the Cyclic Redundancy Check (CRC) code which is transmitted after either the Control Field (for messages with 0 data bytes) or the Data Field, and is used to check the CRC field of incoming messages.

1.6 Error Management Logic

The Error Management Logic is responsible for the fault confinement of the CAN device. Its two counters, the Receive Error Counter (REC) and the Transmit Error Counter (TEC), are incremented and decremented by commands from the Bit Stream Processor. According to the values of the error counters, the CAN controller is set into the states error-active, error-passive or bus-off.

1.7 Bit Timing Logic

The Bit Timing Logic (BTL) monitors the bus line input and handles the bus related bit timing according to the CAN protocol. The BTL synchronizes on a recessive to dominant bus transition at Start of Frame (hard synchronization) and on any further recessive to dominant bus line transition if the CAN controller itself does not transmit a dominant bit (resynchronization). The BTL also provides programmable time segments to compensate for the propagation delay time, phase shifts, and to define the position of the Sample Point within the bit time. The programming of the BTL depends upon the baud rate and external physical delay times.

2.0 CAN MESSAGE FRAMES

The MCP2510 supports Standard Data Frames, Extended Data Frames, and Remote Frames (Standard and Extended) as defined in the CAN 2.0B specification.

2.1 Standard Data Frame

The CAN Standard Data Frame is shown in [Figure 2-1.](#page-8-0) In common with all other frames, the frame begins with a Start Of Frame (SOF) bit, which is of the dominant state, which allows hard synchronization of all nodes.

The SOF is followed by the arbitration field, consisting of 12 bits; the 11-bit ldentifier and the Remote Transmission Request (RTR) bit. The RTR bit is used to distinguish a data frame (RTR bit dominant) from a remote frame (RTR bit recessive).

Following the arbitration field is the control field, consisting of six bits. The first bit of this field is the Identifier Extension (IDE) bit which must be dominant to specify a standard frame. The following bit, Reserved Bit Zero (RB0), is reserved and is defined to be a dominant bit by the can protocol. the remaining four bits of the control field are the Data Length Code (DLC) which specifies the number of bytes of data contained in the message.

After the control field is the data field, which contains any data bytes that are being sent, and is of the length defined by the DLC above (0-8 bytes).

The Cyclic Redundancy Check (CRC) Field follows the data field and is used to detect transmission errors. The CRC Field consists of a 15-bit CRC sequence, followed by the recessive CRC Delimiter bit.

The final field is the two-bit acknowledge field. During the ACK Slot bit, the transmitting node sends out a recessive bit. Any node that has received an error free frame acknowledges the correct reception of the frame by sending back a dominant bit (regardless of whether the node is configured to accept that specific message or not). The recessive acknowledge delimiter completes the acknowledge field and may not be overwritten by a dominant bit.

2.2 Extended Data Frame

In the Extended CAN Data Frame, the SOF bit is followed by the arbitration field which consists of 32 bits, as shown in [Figure 2-2](#page-9-0). The first 11 bits are the most significant bits (Base-lD) of the 29-bit identifier. These 11 bits are followed by the Substitute Remote Request (SRR) bit which is defined to be recessive. The SRR bit is followed by the lDE bit which is recessive to denote an extended CAN frame.

It should be noted that if arbitration remains unresolved after transmission of the first 11 bits of the identifier, and one of the nodes involved in the arbitration is sending a standard CAN frame (11-bit identifier), then the standard CAN frame will win arbitration due to the assertion of a dominant lDE bit. Also, the SRR bit in an extended CAN frame must be recessive to allow the assertion of a dominant RTR bit by a node that is sending a standard CAN remote frame.

The SRR and lDE bits are followed by the remaining 18 bits of the identifier (Extended lD) and the remote transmission request bit.

To enable standard and extended frames to be sent across a shared network, it is necessary to split the 29 bit extended message identifier into 11-bit (most significant) and 18-bit (least significant) sections. This split ensures that the lDE bit can remain at the same bit position in both standard and extended frames.

Following the arbitration field is the six-bit control field. the first two bits of this field are reserved and must be dominant. the remaining four bits of the control field are the Data Length Code (DLC) which specifies the number of data bytes contained in the message.

The remaining portion of the frame (data field, CRC field, acknowledge field, end of frame and lntermission) is constructed in the same way as for a standard data frame (see [Section 2.1](#page-6-1)).

2.3 Remote Frame

Normally, data transmission is performed on an autonomous basis by the data source node (e.g. a sensor sending out a data frame). It is possible, however, for a destination node to request data from the source. To accomplish this, the destination node sends a remote frame with an identifier that matches the identifier of the required data frame. The appropriate data source node will then send a data frame in response to the remote frame request.

There are two differences between a remote frame (shown in [Figure 2-3](#page-10-0)) and a data frame. First, the RTR bit is at the recessive state, and second, there is no data field. In the event of a data frame and a remote frame with the same identifier being transmitted at the same time, the data frame wins arbitration due to the dominant RTR bit following the identifier. In this way, the node that transmitted the remote frame receives the desired data immediately.

2.4 Error Frame

An Error Frame is generated by any node that detects a bus error. An error frame, shown in [Figure 2-4](#page-11-0), consists of two fields, an error flag field followed by an error delimiter field. There are two types of error flag fields. Which type of error flag field is sent depends upon the error status of the node that detects and generates the error flag field.

If an error-active node detects a bus error then the node interrupts transmission of the current message by generating an active error flag. The active error flag is composed of six consecutive dominant bits. This bit

sequence actively violates the bit stuffing rule. All other stations recognize the resulting bit stuffing error and in turn generate error frames themselves, called error echo flags. The error flag field, therefore, consists of between six and twelve consecutive dominant bits (generated by one or more nodes). The error delimiter field completes the error frame. After completion of the error frame, bus activity returns to normal and the interrupted node attempts to resend the aborted message.

If an error-passive node detects a bus error then the node transmits an error-passive flag followed by the error delimiter field. The error-passive flag consists of six consecutive recessive bits, and the error frame for an error-passive node consists of 14 recessive bits. From this, it follows that unless the bus error is detected by the node that is actually transmitting, the transmission of an error frame by an error-passive node will not affect any other node on the network. If the transmitting node generates an error-passive flag then this will cause other nodes to generate error frames due to the resulting bit stuffing violation. After transmission of an error frame, an error-passive node must wait for six consecutive recessive bits on the bus before attempting to rejoin bus communications.

The error delimiter consists of eight recessive bits and allows the bus nodes to restart bus communications cleanly after an error has occurred.

2.5 Overload Frame

An Overload Frame, shown in [Figure 2-5](#page-12-0), has the same format as an active error frame. An overload frame, however can only be generated during an lnterframe space. In this way an overload frame can be differentiated from an error frame (an error frame is sent during the transmission of a message). The overload frame consists of two fields, an overload flag followed by an overload delimiter. The overload flag consists of six dominant bits followed by overload flags generated by other nodes (and, as for an active error flag, giving a maximum of twelve dominant bits). The overload delimiter consists of eight recessive bits. An overload frame can be generated by a node as a result of two conditions. First, the node detects a dominant bit during the interframe space which is an illegal condition. Second, due to internal conditions the node is not yet able to start reception of the next message. A node may generate a maximum of two sequential overload frames to delay the start of the next message.

2.6 Interframe Space

The lnterframe Space separates a preceeding frame (of any type) from a subsequent data or remote frame. The interframe space is composed of at least three recessive bits called the Intermission. This is provided to allow nodes time for internal processing before the start of the next message frame. After the intermission, the bus line remains in the recessive state (bus idle) until the next transmission starts.

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Remote Data Frame with Extended Identifier

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OverloadDelimiter

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Inter-Frame Space or Error Frame

NOTES:

3.0 MESSAGE TRANSMISSION

3.1 Transmit Buffers

The MCP2510 implements three Transmit Buffers. Each of these buffers occupies 14 bytes of SRAM and are mapped into the device memory maps. The first byte, TXBNCTRL, is a control register associated with the message buffer. The information in this register determines the conditions under which the message will be transmitted and indicates the status of the message transmission. (see [Register 3-2\)](#page-17-0). Five bytes are used to hold the standard and extended identifiers and other message arbitration information (see [Register 3-](#page-17-1) [3](#page-17-1) through [Register 3-8\)](#page-19-0). The last eight bytes are for the eight possible data bytes of the message to be transmitted (see [Register 3-8\)](#page-19-0).

For the MCU to have write access to the message buffer, the TXBNCTRL.TXREQ bit must be clear, indicating that the message buffer is clear of any pending message to be transmitted. At a minimum, the TXBN-SIDH, TXBNSIDL, and TXBNDLC registers must be loaded. If data bytes are present in the message, the TXBNDm registers must also be loaded. If the message is to use extended identifiers, the TXBNEIDm registers must also be loaded and the TXBNSIDL.EXIDE bit set.

Prior to sending the message, the MCU must initialize the CANINTE.TXINE bit to enable or disable the generation of an interrupt when the message is sent. The MCU must also initialize the TXBNCTRL.TXP priority bits (see [Section 3.2](#page-14-1)).

3.2 Transmit Priority

Transmit priority is a prioritization, within the MCP2510, of the pending transmittable messages. This is independent from, and not necessarily related to, any prioritization implicit in the message arbitration scheme built into the CAN protocol. Prior to sending the SOF, the priority of all buffers that are queued for transmission is compared. The transmit buffer with the highest priority will be sent first. For example, if transmit buffer 0 has a higher priority setting than transmit buffer 1, buffer 0 will be sent first. If two buffers have the same priority setting, the buffer with the highest buffer number will be sent first. For example, if transmit buffer 1 has the same priority setting as transmit buffer 0, buffer 1 will be sent first. There are four levels of transmit priority. If TXBNC-TRL.TXP<1:0> for a particular message buffer is set to 11, that buffer has the highest possible priority. If TXBNCTRL.TXP<1:0> for a particular message buffer is 00, that buffer has the lowest possible priority.

3.3 Initiating Transmission

To initiate message transmission the TXBNC-TRL.TXREQ bit must be set for each buffer to be transmitted. This can be done by writing to the register via the SPI interface or by setting the TXNRTS pin low for the particular transmit buffer(s) that are to be transmitted. If transmission is initiated via the SPI interface, the TXREQ bit can be set at the same time as the TXP priority bits.

When TXBNCTRL.TXREQ is set, the TXBNCTRL.ABTF, TXBNCTRL.MLOA and TXBNCTRL.TXERR bits will be cleared.

Setting the TXBNCTRL.TXREQ bit does not initiate a message transmission, it merely flags a message buffer as ready for transmission. Transmission will start when the device detects that the bus is available. The device will then begin transmission of the highest priority message that is ready.

When the transmission has completed successfully the TXBNCTRL.TXREQ bit will be cleared, the CAN-INTF.TXNIF bit will be set, and an interrupt will be generated if the CANINTE.TXNIE bit is set.

If the message transmission fails, the TXBNC-TRL.TXREQ will remain set indicating that the message is still pending for transmission and one of the following condition flags will be set. If the message started to transmit but encountered an error condition, the TXBNCTRL. TXERR and the CANINTF.MERRF bits will be set and an interrupt will be generated on the INT pin if the CANINTE.MERRE bit is set. If the message lost arbitration the TXBNCTRL.MLOA bit will be set.

3.4 TXnRTS Pins

The TXNRTS Pins are input pins that can be configured as request-to-send inputs, which provides a secondary means of initiating the transmission of a message from any of the transmit buffers, or as standard digital inputs. Configuration and control of these pins is accomplished using the TXRTSCTRL register (see [Register 3-2](#page-17-0)). The TXRTSCTRL register can only be modified when the MCP2510 is in configuration mode (see [Section 9.0](#page-50-1)). If configured to operate as a request to send pin, the pin is mapped into the respective TXBNCTRL.TXREQ bit for the transmit buffer. The TXREQ bit is latched by the falling edge of the TXNRTS pin. The TXNRTS pins are designed to allow them to be tied directly to the RXNBF pins to automatically initiate a message transmission when the RXNBF pin goes low. The TXNRTS pins have internal pullup resistors of 100 kΩ (nominal).

3.5 Aborting Transmission

The MCU can request to abort a message in a specific message buffer by clearing the associated TXBnC-TRL.TXREQ bit. Also, all pending messages can be requested to be aborted by setting the CAN-CTRL.ABAT bit. If the CANCTRL.ABAT bit is set to abort all pending messages, the user MUST reset this bit (typically after the user verifies that all TXREQ bits have been cleared) to continue trasmit messages. The CANCTRL.ABTF flag will only be set if the abort was requested via the CANCTRL.ABAT bit. Aborting a message by resetting the TXREQ bit does cause the ATBF bit to be set.

Only messages that have not already begun to be transmitted can be aborted. Once a message has begun transmission, it will not be possible for the user to reset the TXBnCTRL.TXREQ bit. After transmission

of a message has begun, if an error occurs on the bus or if the message loses arbitration, the message will be retransmitted regardless of a request to abort.

REGISTER 3-1: TXBNCTRL Transmit Buffer N Control Register (ADDRESS: 30h, 40h, 50h)

	$U-0$	$R - 0$	$R - 0$	$R - 0$	$R/W-0$	$U-0$	$R/W-0$	$R/W-0$
		ABTF	MLOA	TXERR	TXREQ		TXP1	TXP ₀
	bit 7							bit 0
bit 7	Unimplemented: Read as '0'							
bit 6		ABTF: Message Aborted Flag						
		$1 =$ Message was aborted	$0 =$ Message completed transmission successfully					
bit 5		MLOA: Message Lost Arbitration						
	$1 =$ Message lost arbitration while being sent $0 =$ Message did not lose arbitration while being sent							
bit 4			TXERR: Transmission Error Detected					
					$1 =$ A bus error occurred while the message was being transmitted $0 =$ No bus error occurred while the message was being transmitted			
bit 3			TXREQ: Message Transmit Request					
	$1 =$ Buffer is currently pending transmission (MCU sets this bit to request message be transmitted - bit is automatically cleared when the message is sent) $0 =$ Buffer is not currently pending transmission (MCU can clear this bit to request a message abort)							
bit 2	Unimplemented: Read as '0'							
bit 1-0	TXP<1:0>: Transmit Buffer Priority							
	$11 =$ $10 =$ $11 =$ $00 =$	Highest Message Priority Lowest Message Priority	High Intermediate Message Priority Low Intermediate Message Priority					
	Legend:							

REGISTER 3-2: TXRTSCTRL - TXNRTS PIN CONTROL AND STATUS REGISTER (ADDRESS: 0Dh)

REGISTER 3-3: TXBNSIDH - TRANSMIT BUFFER N STANDARD IDENTIFIER HIGH (ADDRESS: 31h, 41h, 51h)

bit 7-0

REGISTER 3-5: TXBNEID8 - TRANSMIT BUFFER N EXTENDED IDENTIFIER HIGH (ADDRESS: 33h, 43h, 53h)

bit 7-0 **EID<15:8>**: Extended Identifier Bits <15:8>

REGISTER 3-6: TXBNEID0 - TRANSMIT BUFFER N EXTENDED IDENTIFIER LOW (ADDRESS: 34h, 44h, 54h)

bit 7-0 **EID<7:0>**: Extended Identifier Bits <7:0>

REGISTER 3-8: TXBNDM - Transmit Buffer N Data Field Byte m (ADDRESS: 36h-3Dh, 46h-4Dh, 56h-5Dh)

bit 7-0 **TXBNDM7:TXBNDM0**: Transmit Buffer N Data Field Byte m

4.0 MESSAGE RECEPTION

4.1 Receive Message Buffering

The MCP2510 includes two full receive buffers with multiple acceptance filters for each. There is also a separate Message Assembly Buffer (MAB) which acts as a third receive buffer (see [Figure 4-1](#page-21-0)).

4.2 Receive Buffers

Of the three Receive Buffers, the MAB is always committed to receiving the next message from the bus. The remaining two receive buffers are called RXB0 and RXB1 and can receive a complete message from the protocol engine. The MCU can access one buffer while the other buffer is available for message reception or holding a previously received message.

The MAB assembles all messages received. These messages will be transferred to the RXBN buffers (See [Register 4-4](#page-25-0) to [Register 4-9\)](#page-27-0) only if the acceptance filter criteria are met.

Note: The entire contents of the MAB is moved into the receive buffer once a message is accepted. This means that regardless of the type of identifier (standard or extended) and the number of data bytes received, the entire receive buffer is overwritten with the MAB contents. Therefore the contents of all registers in the buffer must be assumed to have been modified when any message is received.

When a message is moved into either of the receive buffers the appropriate CANINTF.RXNIF bit is set. This bit must be cleared by the MCU, when it has completed processing the message in the buffer, in order to allow a new message to be received into the buffer. This bit provides a positive lockout to ensure that the MCU has finished with the message before the MCP2510 attempts to load a new message into the receive buffer. If the CANINTE.RXNIE bit is set an interrupt will be generated on the $\overline{\text{INT}}$ pin to indicate that a valid message has been received.

4.3 Receive Priority

RXB0 is the higher priority buffer and has two message acceptance filters associated with it. RXB1 is the lower priority buffer and has four acceptance filters associated with it. The lower number of acceptance filters makes the match on RXB0 more restrictive and implies a higher priority for that buffer. Additionally, the RXB0CTRL register can be configured such that if RXB0 contains a valid message, and another valid message is received, an overflow error will not occur and the new message will be moved into RXB1 regardless of the acceptance criteria of RXB1. There are also two programmable acceptance filter masks available, one for each receive buffer (see [Section 4.5](#page-28-0)).

When a message is received, bits <3:0> of the RXBNC-TRL Register will indicate the acceptance filter number that enabled reception, and whether the received message is a remote transfer request.

The RXBNCTRL.RXM bits set special receive modes. Normally, these bits are set to 00 to enable reception of all valid messages as determined by the appropriate acceptance filters. In this case, the determination of whether or not to receive standard or extended messages is determined by the RFXNSIDL.EXIDE bit in the acceptance filter register. If the RXBNCTRL.RXM bits are set to 01 or 10, the receiver will accept only messages with standard or extended identifiers respectively. If an acceptance filter has the RFXNSIDL.EXIDE bit set such that it does not correspond with the RXBNCTRL.RXM mode, that acceptance filter is rendered useless. These two modes of RXBNCTRL.RXM bits can be used in systems where it is known that only standard or extended messages will be on the bus. If the RXBNCTRL.RXM bits are set to 11, the buffer will receive all messages regardless of the values of the acceptance filters. Also, if a message has an error before the end of frame, that portion of the message assembled in the MAB before the error frame will be loaded into the buffer. This mode has some value in debugging a CAN system and would not be used in an actual system environment.

4.4 RX0BF and RX1BF Pins

In addition to the $\overline{\text{INT}}$ pin which provides an interrupt signal to the MCU for many different conditions, the receive buffer full pins (RX0BF and RX1BF) can be used to indicate that a valid message has been loaded into RXB0 or RXB1, respectively.

The RXBNBF full pins can be configured to act as buffer full interrupt pins or as standard digital outputs. Configuration and status of these pins is available via the BFPCTRL register ([Register 4-3\)](#page-25-1). When set to operate in interrupt mode (by setting BFPCTRL.BxBFE and BFPCTRL.BxBFM bits to a 1), these pins are active low and are mapped to the CANINTF.RXNIF bit for each receive buffer. When this bit goes high for one of the receive buffers, indicating that a valid message has been loaded into the buffer, the corresponding RXNBF pin will go low. When the CANINTF.RXNIF bit is cleared by the MCU, then the corresponding interrupt pin will go to the logic high state until the next message is loaded into the receive buffer.

When used as digital outputs, the BFPCTRL.BxBFM bits must be cleared to a '0' and BFPCTRL.BxBFE bits must be set to a '1' for the associated buffer. In this mode the state of the pin is controlled by the BFPC-TRL.BxBFS bits. Writting a '1' to the BxBFS bit will cause a high level to be driven on the assicated buffer full pin, and a '0' will cause the pin to drive low. When using the pins in this mode the state of the pin should be modified only by using the Bit Modify SPI command to prevent glitches from occuring on either of the buffer full pins.

REGISTER 4-1: RXB0CTRL - RECEIVE BUFFER 0 CONTROL REGISTER (ADDRESS: 60h)

REGISTER 4-2: RXB1CTRL - RECEIVE BUFFER 1 CONTROL REGISTER (ADDRESS: 70h)

REGISTER 4-3: BFPCTRL - RXNBF PIN CONTROL AND STATUS REGISTER (ADDRESS: 0Ch)

REGISTER 4-4: RXBNSIDH - RECEIVE BUFFER N STANDARD IDENTIFIER HIGH (ADDRESS: 61h, 71h)

bit 7-0 **SID<10:3>**: Standard Identifier Bits <10:3>

These bits contain the eight most significant bits of the Standard Identifier for the received message

REGISTER 4-5: RXBNSIDL - RECEIVE BUFFER N STANDARD IDENTIFIER LOW (ADDRESS: 62h, 72h)

REGISTER 4-6: RXBNEID8 - RECEIVE BUFFER N EXTENDED IDENTIFIER MID (ADDRESS: 63h, 73h)

bit 7-0 **EID<15:8>**: Extended Identifier Bits <15:8> These bits hold bits 15 through 8 of the Extended Identifier for the received message

REGISTER 4-7: RXBNEID0 - RECEIVE BUFFER N EXTENDED IDENTIFIER LOW (ADDRESS: 64h, 74h)

bit 7-0 **EID<7:0>**: Extended Identifier Bits <7:0>

These bits hold the least significant eight bits of the Extended Identifier for the received message

REGISTER 4-8: RXBNDLC - RECEIVE BUFFER N DATA LENGTH CODE (ADDRESS: 65h, 75h)

bit 7 **Unimplemented:** Reads as '0'

bit 6 **RTR**: Extended Frame Remote Transmission Request Bit (valid only when RXBnSIDL.IDE = 1) 1 = Extended Frame Remote Transmit Request Received 0 = Extended Data Frame Received

- bit 5 **RB1:** Reserved Bit 1
- bit 4 **RB0:** Reserved Bit 0
- bit 3-0 **DLC<3:0>**: Data Length Code Indicates number of data bytes that were received

REGISTER 4-9: RXBNDM - RECEIVE BUFFER N DATA FIELD BYTE M (ADDRESS: 66h-6Dh, 76h-7Dh)

bit 7-0 **RBNDm7:RBNDm0**: Receive Buffer N Data Field Byte m

Eight bytes containing the data bytes for the received message

4.5 Message Acceptance Filters and Masks

The Message Acceptance Filters And Masks are used to determine if a message in the message assembly buffer should be loaded into either of the receive buffers (see [Figure 4-3\)](#page-29-0). Once a valid message has been received into the MAB, the identifier fields of the message are compared to the filter values. If there is a match, that message will be loaded into the appropriate receive buffer. The filter masks (see [Register 4-10](#page-29-1) through [Register 4-17\)](#page-32-0) are used to determine which bits in the identifier are examined with the filters. A truth table is shown below in [Table 4-1](#page-28-1) that indicates how each bit in the identifier is compared to the masks and filters to determine if a the message should be loaded into a receive buffer. The mask essentially determines which bits to apply the acceptance filters to. If any mask bit is set to a zero, then that bit will automatically be accepted regardless of the filter bit.

TABLE 4-1: FILTER/MASK TRUTH TABLE

Note: X = don't care

As shown in the Receive Buffers Block Diagram [\(Figure 4-1\)](#page-21-0), acceptance filters RXF0 and RXF1, and filter mask RXM0 are associated with RXB0. Filters RXF2, RXF3, RXF4, and RXF5 and mask RXM1 are associated with RXB1. When a filter matches and a message is loaded into the receive buffer, the filter number that enabled the message reception is loaded into the RXBNCTRL register FILHIT bit(s). For RXB1 the RXB1CTRL register contains the FILHIT<2:0> bits. They are coded as follows:

- 101 = Acceptance Filter 5 (RXF5)
- 100 = Acceptance Filter 4 (RXF4)
- 011 = Acceptance Filter 3 (RXF3)
- 010 = Acceptance Filter 2 (RXF2)
- 001 = Acceptance Filter 1 (RXF1)
- 000 = Acceptance Filter 0 (RXF0)

RXB0CTRL contains two copies of the BUKT bit and the FILHIT<0> bit.

The coding of the BUKT bit enables these three bits to be used similarly to the RXB1CTRL.FILHIT bits and to distinguish a hit on filter RXF0 and RXF1 in either RXB0 or after a roll over into RXB1.

- -111 = Acceptance Filter 1 (RXF1)
- 110 = Acceptance Filter 0 (RXF0)
- -001 = Acceptance Filter 1 (RXF1)
- 000 = Acceptance Filter 0

If the BUKT bit is clear, there are six codes corresponding to the six filters. If the BUKT bit is set, there are six codes corresponding to the six filters plus two additional codes corresponding to RXF0 and RXF1 filters that roll over into RXB1.

If more than one acceptance filter matches, the FILHIT bits will encode the binary value of the lowest numbered filter that matched. In other words, if filter RXF2 and filter RXF4 match, FILHIT will be loaded with the value for RXF2. This essentially prioritizes the acceptance filters with a lower number filter having higher priority. Messages are compared to filters in ascending order of filter number.

The mask and filter registers can only be modified when the MCP2510 is in configuration mode (see [Section 9.0](#page-50-1)).

FIGURE 4-3: MESSAGE ACCEPTANCE MASK AND FILTER OPERATION

REGISTER 4-10: RXFNSIDH - ACCEPTANCE FILTER N STANDARD IDENTIFIER HIGH (ADDRESS: 00h, 04h, 08h, 10h, 14h, 18h)

bit 7-0 **SID<10:3>**: Standard Identifier Filter Bits <10:3>

These bits hold the filter bits to be applied to bits <10:3> of the Standard Identifier portion of a received message

REGISTER 4-11: RXFNSIDL - ACCEPTANCE FILTER N STANDARD IDENTIFIER LOW (ADDRESS: 01h, 05h, 09h, 11h, 15h, 19h)

REGISTER 4-12: RXFNEID8 - ACCEPTANCE FILTER N EXTENDED IDENTIFIER HIGH (ADDRESS: 02h, 06h, 0Ah, 12h, 16h, 1Ah)

 $-n =$ Value at POR $'1' =$ Bit is set $'0' =$ Bit is cleared $x =$ Bit is unknown

bit 7-0 **EID<15:8>**: Extended Identifier Bits <15:8> These bits hold the filter bits to be applied to bits <15:8> of the Extended Identifier portion of a received message

REGISTER 4-13: RXFNEID0 - ACCEPTANCE FILTER N EXTENDED IDENTIFIER LOW (ADDRESS: 03h, 07h, 0Bh, 13h, 17h, 1Bh)

bit 7-0 **EID<7:0>**: Extended Identifier Bits <7:0>

These bits hold the filter bits to be applied to the bits <7:0> of the Extended Identifier portion of a received message

REGISTER 4-14: RXMNSIDH - ACCEPTANCE FILTER MASK N STANDARD IDENTIFIER HIGH (ADDRESS: 20h, 24h)

bit 7-0 **SID<10:3>**: Standard Identifier Mask Bits <10:3>

These bits hold the mask bits to be applied to bits <10:3> of the Standard Identifier portion of a received message

REGISTER 4-15: RXMNSIDL - ACCEPTANCE FILTER MASK N STANDARD IDENTIFIER LOW (ADDRESS: 21h, 25h)

- bit 7-5 **SID<2:0>**: Standard Identifier Mask Bits <2:0> These bits hold the mask bits to be applied to bits<2:0> of the Standard Identifier portion of a received message
- bit 4-2 **Unimplemented**: Reads as '0'
- bit 1-0 **EID<17:16>**: Extended Identifier Mask Bits <17:16> These bits hold the mask bits to be applied to bits <17:16> of the Extended Identifier portion of a received message

REGISTER 4-16: RXMNEID8 - ACCEPTANCE FILTER MASK N EXTENDED IDENTIFIER HIGH (ADDRESS: 22h, 26h)

bit 7-0 **EID<15:8>**: Extended Identifier Bits <15:8>

These bits hold the filter bits to be applied to bits <15:8> of the Extended Identifier portion of a received message

REGISTER 4-17: RXMNEID0 - ACCEPTANCE FILTER MASK N EXTENDED IDENTIFIER LOW (ADDRESS: 23h, 27h)

bit 7-0 **EID<7:0>**: Extended Identifier Mask Bits <7:0>

These bits hold the mask bits to be applied to the bits <7:0> of the Extended Identifier portion of a received message

NOTES:

5.0 BIT TIMING

All nodes on a given CAN bus must have the same nominal bit rate. The CAN protocol uses Non Return to Zero (NRZ) coding which does not encode a clock within the data stream. Therefore, the receive clock must be recovered by the receiving nodes and synchronized to the transmitters clock.

As oscillators and transmission time may vary from node to node, the receiver must have some type of Phase Lock Loop (PLL) synchronized to data transmission edges to synchronize and maintain the receiver clock. Since the data is NRZ coded, it is necessary to include bit stuffing to ensure that an edge occurs at least every six bit times, to maintain the Digital Phase Lock Loop (DPLL) synchronization.

The bit timing of the MCP2510 is implemented using a DPLL that is configured to synchronize to the incoming data, and provide the nominal timing for the transmitted data. The DPLL breaks each bit time into multiple segments made up of minimal periods of time called the time quanta (TQ).

Bus timing functions executed within the bit time frame, such as synchronization to the local oscillator, network transmission delay compensation, and sample point positioning, are defined by the programmable bit timing logic of the DPLL.

All devices on the CAN bus must use the same bit rate. However, all devices are not required to have the same master oscillator clock frequency. For the different clock frequencies of the individual devices, the bit rate has to be adjusted by appropriately setting the baud rate prescaler and number of time quanta in each segment.

FIGURE 5-1: BIT TIME PARTITIONING

Nominal Bit Time is defined as:

TBIT = 1 / NOMlNAL BlT RATE

The nominal bit time can be thought of as being divided into separate non-overlapping time segments. These segments are shown in [Figure 5-1.](#page-34-1)

- Synchronization Segment (Sync_Seg)
- Propagation Time Segment (Prop_Seg)
- Phase Buffer Segment 1 (Phase_Seg1)
- Phase Buffer Segment 2 [Phase_Seg2)

Nominal Bit Time = TQ * (Sync_Seg + Prop_Seg + Phase_Seg1 + Phase_Seg2)

The time segments and also the nominal bit time are made up of integer units of time called time quanta or To (see [Figure 5-1](#page-34-1)). By definition, the nominal bit time is programmable from a minimum of 8 TQ to a maximum of 25 TQ. Also, by definition the minimum nominal bit time is 1 µs, corresponding to a maximum 1 Mb/s rate.

5.1 Time Quanta

The Time Quanta (TQ) is a fixed unit of time derived from the oscillator period. There is a programmable baud-rate prescaler, with integral values ranging from 1 to 64, in addition to a fixed divide by two for clock generation.

Time quanta is defined as:

 $T_{Q} = 2*($ Baud Rate + 1)^{*} T_{QSC}

where Baud Rate is the binary value represented by CNF1.BRP<5:0>

For some examples:

If Fosc = 16 MHz, BRP $<5:0>$ = 00h, and Nominal Bit $Time = 8$ T Q :

then $To = 125$ nsec and Nominal Bit Rate = 1 Mb/s

If FOSC = 20 MHz, $BRP<5:0>$ = 01h, and Nominal Bit $Time = 8$ T Q :

then $TQ = 200$ nsec and Nominal Bit Rate = 625 Kb/s If Fosc = 25 MHz, BRP< $5:0$ > = $3Fh$, and Nominal Bit $Time = 25$ T Q ;

then $TQ = 5.12$ µsec and Nominal Bit Rate = 7.8 Kb/s

The frequencies of the oscillators in the different nodes must be coordinated in order to provide a system-wide specified nominal bit time. This means that all oscillators must have a TOSC that is a integral divisor of TQ. It should also be noted that although the number of TQ is programmable from 4 to 25, the usable minimum is 6 TQ. Attempting to a bit time of less than 6 TQ in length is not guaranteed to operate correctly

5.2 Synchronization Segment

This part of the bit time is used to synchronize the various CAN nodes on the bus. The edge of the input signal is expected to occur during the sync segment. The duration is 1 TQ.

5.3 Propagation Segment

This part of the bit time is used to compensate for physical delay times within the network. These delay times consist of the signal propagation time on the bus line and the internal delay time of the nodes. The delay is calculated as being the round trip time from transmitter to receiver (twice the signal's propagation time on the bus line), the input comparator delay, and the output driver delay. The length of the Propagation Segment can be programmed from 1 TQ to 8 TQ by setting the PRSEG2:PRSEG0 bits of the CNF2 register [\(Register 5-2\)](#page-39-0).

The total delay is calculated from the following individual delays:

- 2 * physical bus end to end delay; TBUS
- 2^{*} input comparator delay; TCOMP (depends on application circuit)
- 2 * output driver delay; TDRIVE (depends on application circuit)
- 1^{*} input to output of CAN controller; TCAN (maximum defined as 1 T Q + delay ns)
- TPROPOGATION = $2 *$ (TBUS + TCOMP + TDRIVE) + TCAN
- Prop_Seg = TPROPOGATION / TQ

5.4 Phase Buffer Segments

The Phase Buffer Segments are used to optimally locate the sampling point of the received bit within the nominal bit time. The sampling point occurs between phase segment 1 and phase segment 2. These segments can be lengthened or shortened by the resynchronization process (see [Section 5.7.2](#page-36-0)). Thus, the variation of the values of the phase buffer segments represent the DPLL functionality. The end of phase segment 1 determines the sampling point within a bit time. phase segment 1 is programmable from 1 To to 8 TQ in duration. Phase segment 2 provides delay before the next transmitted data transition and is also programmable from 1 TQ to 8 TQ in duration (however due to IPT requirements the actual minimum length of phase segment 2 is 2 TQ - see [Section 5.6](#page-35-0) below), or it may be defined to be equal to the greater of phase segment 1 or the Information Processing Time (IPT). (see [Section 5.6](#page-35-0)).

5.5 Sample Point

The Sample Point is the point of time at which the bus level is read and value of the received bit is determined. The Sampling point occurs at the end of phase segment 1. If the bit timing is slow and contains many TQ, it is possible to specify multiple sampling of the bus line at the sample point. The value of the received bit is determined to be the value of the majority decision of three values. The three samples are taken at the sample point, and twice before with a time of TQ/2 between each sample.

5.6 Information Processing Time

The Information Processing Time (IPT) is the time segment, starting at the sample point, that is reserved for calculation of the subsequent bit level. The CAN specification defines this time to be less than or equal to 2 TQ. The MCP2510 defines this time to be 2 TQ. Thus, phase segment 2 must be at least 2 TQ long.
5.7 Synchronization

To compensate for phase shifts between the oscillator frequencies of each of the nodes on the bus, each CAN controller must be able to synchronize to the relevant signal edge of the incoming signal. Synchronization is the process by which the DPLL function is implemented. When an edge in the transmitted data is detected, the logic will compare the location of the edge to the expected time (Sync Seg). The circuit will then adjust the values of phase segment 1 and phase segment 2 as necessary. There are two mechanisms used for synchronization.

5.7.1 HARD SYNCHRONIZATION

Hard Synchronization is only done when there is a recessive to dominant edge during a BUS IDLE condition, indicating the start of a message. After hard synchronization, the bit time counters are restarted with Sync Seg. Hard synchronization forces the edge which has occurred to lie within the synchronization segment of the restarted bit time. Due to the rules of synchronization, if a hard synchronization occurs there will not be a resynchronization within that bit time.

5.7.2 RESYNCHRONIZATION

As a result of Resynchronization, phase segment 1 may be lengthened or phase segment 2 may be shortened. The amount of lengthening or shortening of the phase buffer segments has an upper bound given by the Synchronization Jump Width (SJW). The value of the SJW will be added to phase segment 1 (see [Figure 5-2\)](#page-36-0) or subtracted from phase segment 2 (see [Figure 5-3\)](#page-37-0). The SJW represents the loop filtering of the DPLL. The SJW is programmable between 1 TQ and 4 TQ.

Clocking information will only be derived from recessive to dominant transitions. The property that only a fixed maximum number of successive bits have the same value ensures resynchronization to the bit stream during a frame.

The phase error of an edge is given by the position of the edge relative to Sync Seg, measured in TQ. The phase error is defined in magnitude of TQ as follows:

- \cdot e = 0 if the edge lies within SYNCESEG
- \cdot e > 0 if the edge lies before the SAMPLE POINT
- e < 0 if the edge lies after the SAMPLE POINT of the previous bit

If the magnitude of the phase error is less than or equal to the programmed value of the synchronization jump width, the effect of a resynchronization is the same as that of a hard synchronization.

If the magnitude of the phase error is larger than the synchronization jump width, and if the phase error is positive, then phase segment 1 is lengthened by an amount equal to the synchronization jump width.

If the magnitude of the phase error is larger than the resynchronization jump width, and if the phase error is negative, then phase segment 2 is shortened by an amount equal to the synchronization jump width.

5.7.3 SYNCHRONIZATION RULES

- Only one synchronization within one bit time is allowed
- An edge will be used for synchronization only if the value detected at the previous sample point (previously read bus value) differs from the bus value immediately after the edge
- All other recessive to dominant edges fulfilling rules 1 and 2 will be used for resynchronization with the exception that a node transmitting a dominant bit will not perform a resynchronization as a result of a recessive to dominant edge with a positive phase error

FIGURE 5-2: LENGTHENING A BIT PERIOD

FIGURE 5-3: SHORTENING A BIT PERIOD

5.8 Programming Time Segments

Some requirements for programming of the time segments:

- Prop Seg + Phase Seg 1 >= Phase Seg 2
- Prop Seg + Phase Seg 1 >= TDELAY
- Phase Seg 2 > Sync Jump Width

For example, assuming that a 125 kHz CAN baud rate with FOSC = 20 MHz is desired:

Tosc = 50 nsec, choose BRP<5:0> = 04h, then $TQ =$ 500 nsec. To obtain 125 kHz, the bit time must be 16 TQ.

Typically, the sampling of the bit should take place at about 60-70% of the bit time, depending on the system parameters. Also, typically, the TDELAY is 1-2 TQ.

Sync Seg = 1 TQ; Prop Seg = 2 TQ; So setting Phase Seg $1 = 7$ TQ would place the sample at 10 TQ after the transition. This would leave 6 TQ for Phase Seg 2.

Since Phase Seg 2 is 6, by the rules, SJW could be the maximum of 4 TQ. However, normally a large SJW is only necessary when the clock generation of the different nodes is inaccurate or unstable, such as using ceramic resonators. So an SJW of 1 is typically enough.

5.9 Oscillator Tolerance

The bit timing requirements allow ceramic resonators to be used in applications with transmission rates of up to 125 kbit/sec, as a rule of thumb. For the full bus speed range of the CAN protocol, a quartz oscillator is required. A maximum node-to-node oscillator variation of 1.7% is allowed.

5.10 Bit Timing Configuration Registers

The configuration registers (CNF1, CNF2, CNF3) control the bit timing for the CAN bus interface. These registers can only be modified when the MCP2510 is in configuration mode (see [Section 9.0\)](#page-50-0).

5.10.1 CNF1

The BRP<5:0> bits control the baud rate prescaler. These bits set the length of TQ relative to the OSC1 input frequency, with the minimum length of TQ being 2 OSC1 clock cycles in length (when BRP<5:0> are set to 000000). The SJW<1:0> bits select the synchronization jump width in terms of number of TQ's.

5.10.2 CNF2

The PRSEG<2:0> bits set the length, in TQ's, of the propagation segment. The PHSEG1<2:0> bits set the length, in TQ's, of phase segment 1. The SAM bit controls how many times the RXCAN pin is sampled. Setting this bit to a '1' causes the bus to be sampled three times; twice at TQ/2 before the sample point, and once at the normal sample point (which is at the end of phase segment 1). The value of the bus is determined to be the value read during at least two of the samples. If the SAM bit is set to a '0' then the RXCAN pin is sampled only once at the sample point. The BTLMODE bit controls how the length of phase segment 2 is determined. If this bit is set to a '1' then the length of phase segment 2 is determined by the PHSEG2<2:0> bits of CNF3 (see [Section 5.10.3](#page-38-0)). If the BTLMODE bit is set to a '0' then the length of phase segment 2 is the greater of phase segment 1 and the information processing time (which is fixed at 2 TQ for the MCP2510).

5.10.3 CNF3

The PHSEG2<2:0> bits set the length, in TQ's, of Phase Segment 2, if the CNF2.BTLMODE bit is set to a '1'. If the BTLMODE bit is set to a '0' then the PHSEG2<2:0> bits have no effect.

REGISTER 5-1: CNF1 - CONFIGURATION REGISTER1 (ADDRESS: 2Ah)

 $-n =$ Value at POR $i'1' =$ Bit is set $i'0' =$ Bit is cleared $x =$ Bit is unknown

REGISTER 5-2: CNF2 - CONFIGURATION REGISTER2 (ADDRESS: 29h)

REGISTER 5-3: CNF3 - CONFIGURATION REGISTER 3 (ADDRESS: 28h)

6.0 ERROR DETECTION

The CAN protocol provides sophisticated error detection mechanisms. The following errors can be detected.

6.1 CRC Error

With the Cyclic Redundancy Check (CRC), the transmitter calculates special check bits for the bit sequence from the start of a frame until the end of the data field. This CRC sequence is transmitted in the CRC Field. The receiving node also calculates the CRC sequence using the same formula and performs a comparison to the received sequence. If a mismatch is detected, a CRC error has occurred and an error frame is generated. The message is repeated.

6.2 Acknowledge Error

In the acknowledge field of a message, the transmitter checks if the acknowledge slot (which has sent out as a recessive bit) contains a dominant bit. If not, no other node has received the frame correctly. An acknowledge error has occurred; an error frame is generated; and the message will have to be repeated.

6.3 Form Error

lf a node detects a dominant bit in one of the four segments including end of frame, interframe space, acknowledge delimiter or CRC delimiter; then a form error has occurred and an error frame is generated. The message is repeated.

6.4 Bit Error

A Bit Error occurs if a transmitter sends a dominant bit and detects a recessive bit or if it sends a recessive bit and detects a dominant bit when monitoring the actual bus level and comparing it to the just transmitted bit. In the case where the transmitter sends a recessive bit and a dominant bit is detected during the arbitration field and the acknowledge slot, no bit error is generated because normal arbitration is occurring.

6.5 Stuff Error

lf, between the start of frame and the CRC delimiter, six consecutive bits with the same polarity are detected, the bit stuffing rule has been violated. A stuff error occurs and an error frame is generated. The message is repeated.

6.6 Error States

Detected errors are made public to all other nodes via error frames. The transmission of the erroneous message is aborted and the frame is repeated as soon as possible. Furthermore, each CAN node is in one of the three error states "error-active", "error-passive" or "busoff" according to the value of the internal error counters. The error-active state is the usual state where the bus node can transmit messages and active error frames (made of dominant bits) without any restrictions. In the error-passive state, messages and passive error frames (made of recessive bits) may be transmitted. The bus-off state makes it temporarily impossible for the station to participate in the bus communication. During this state, messages can neither be received nor transmitted.

6.7 Error Modes and Error Counters

The MCP2510 contains two error counters: the Receive Error Counter (REC) (see [Register 6-2](#page-41-0)), and the Transmit Error Counter (TEC) (see [Register 6-1\)](#page-41-1). The values of both counters can be read by the MCU. These counters are incremented or decremented in accordance with the CAN bus specification.

The MCP2510 is error-active if both error counters are below the error-passive limit of 128. It is error-passive if at least one of the error counters equals or exceeds 128. It goes to bus-off if the transmit error counter equals or exceeds the bus-off limit of 256. The device remains in this state, until the bus-off recovery sequence is received. The bus-off recovery sequence consists of 128 occurrences of 11 consecutive recessive bits (see [Figure 6-1\)](#page-41-2). Note that the MCP2510, after going bus-off, will recover back to error-active, without any intervention by the MCU, if the bus remains idle for 128 X 11 bit times. If this is not desired, the error interrupt service routine should address this. The current error mode of the MCP2510 can be read by the MCU via the EFLG register [\(Register 6-3](#page-42-0)).

Additionally, there is an error state warning flag bit, EFLG:EWARN, which is set if at least one of the error counters equals or exceeds the error warning limit of 96. EWARN is reset if both error counters are less than the error warning limit.

FIGURE 6-1: ERROR MODES STATE DIAGRAM

REGISTER 6-1: TEC - TRANSMITTER ERROR COUNTER (ADDRESS: 1Ch)

bit 7-0 **TEC<7:0>**: Transmit Error Count

REGISTER 6-2: REC - RECEIVER ERROR COUNTER (ADDRESS: 1Dh)

bit 7-0 **REC<7:0>**: Receive Error Count

REGISTER 6-3: EFLG - ERROR FLAG REGISTER (ADDRESS: 2Dh)

NOTES:

7.0 INTERRUPTS

The device has eight sources of interrupts. The CAN-INTE register contains the individual interrupt enable bits for each interrupt source. The CANINTF register contains the corresponding interrupt flag bit for each interrupt source. When an interrupt occurs the INT pin is driven low by the MCP2510 and will remain low until the Interrupt is cleared by the MCU. An Interrupt can not be cleared if the respective condition still prevails.

It is recommended that the bit modify command be used to reset flag bits in the CANINTF register rather than normal write operations. This is to prevent unintentionally changing a flag that changes during the write command, potentially causing an interrupt to be missed.

It should be noted that the CANINTF flags are read/ write and an Interrupt can be generated by the MCU setting any of these bits, provided the associated CAN-INTE bit is also set.

7.1 Interrupt Code Bits

The source of a pending interrupt is indicated in the CANSTAT.ICOD (interrupt code) bits as indicated in [Register 9-2](#page-52-0). In the event that multiple interrupts occur, the INT will remain low until all interrupts have been reset by the MCU, and the CANSTAT.ICOD bits will reflect the code for the highest priority interrupt that is currently pending. Interrupts are internally prioritized such that the lower the ICOD value the higher the interrupt priority. Once the highest priority interrupt condition has been cleared, the code for the next highest priority interrupt that is pending (if any) will be reflected by the ICOD bits (see [Table 7-1\)](#page-44-0). Note that only those interrupt sources that have their associated CANINTE enable bit set will be reflected in the ICOD bits.

7.2 Transmit Interrupt

When the Transmit Interrupt is enabled (CAN- $INTER$ NIE = 1) an Interrupt will be generated on the INT pin when the associated transmit buffer becomes empty and is ready to be loaded with a new message. The CANINTF.TXNIF bit will be set to indicate the source of the interrupt. The interrupt is cleared by the MCU resetting the TXNIF bit to a '0'.

7.3 Receive Interrupt

When the Receive Interrupt is enabled (CAN- $INTERXNIE = 1)$ an interrupt will be generated on the INT pin when a message has been successfully received and loaded into the associated receive buffer. This interrupt is activated immediately after receiving the EOF field. The CANINTF.RXNIF bit will be set to indicate the source of the interrupt. The interrupt is cleared by the MCU resetting the RXNIF bit to a '0'.

7.4 Message Error Interrupt

When an error occurs during transmission or reception of a message the message error flag (CAN-INTF.MERRF) will be set and, if the CANINTE.MERRE bit is set, an interrupt will be generated on the INT pin. This is intended to be used to facilitate baud rate determination when used in conjunction with listen-only mode.

7.5 Bus Activity Wakeup Interrupt

When the MCP2510 is in sleep mode and the bus activity wakeup interrupt is enabled (CANINTE.WAKIE = 1), an interrupt will be generated on the $\overline{\text{INT}}$ pin, and the CANINTF.WAKIF bit will be set when activity is detected on the CAN bus. This interrupt causes the MCP2510 to exit sleep mode. The interrupt is reset by the MCU clearing the WAKIF bit.

7.6 Error Interrupt

When the error interrupt is enabled (CANINTE.ERRIE $=$ 1) an interrupt is generated on the $\overline{\text{INT}}$ pin if an overflow condition occurs or if the error state of transmitter or receiver has changed. The Error Flag Register (EFLG) will indicate one of the following conditions.

7.6.1 RECEIVER OVERFLOW

An overflow condition occurs when the MAB has assembled a valid received message (the message meets the criteria of the acceptance filters) and the receive buffer associated with the filter is not available for loading of a new message. The associated EFLG.RXNOVR bit will be set to indicate the overflow condition. This bit must be cleared by the MCU.

7.6.2 RECEIVER WARNING

The receive error counter has reached the MCU warning limit of 96.

7.6.3 TRANSMITTER WARNING

The transmit error counter has reached the MCU warning limit of 96.

7.6.4 RECEIVER ERROR-PASSIVE

The receive error counter has exceeded the error- passive limit of 127 and the device has gone to error- passive state.

7.6.5 TRANSMITTER ERROR-PASSIVE

The transmit error counter has exceeded the errorpassive limit of 127 and the device has gone to errorpassive state.

7.6.6 BUS-OFF

The transmit error counter has exceeded 255 and the device has gone to bus-off state.

7.7 Interrupt Acknowledge

Interrupts are directly associated with one or more status flags in the CANINTF register. Interrupts are pending as long as one of the flags is set. Once an interrupt flag is set by the device, the flag can not be reset by the MCU until the interrupt condition is removed.

 $-n =$ Value at POR $'1' = B$ it is set $'0' = B$ it is cleared $x = B$ it is unknown

REGISTER 7-2: CANINTF - INTERRUPT FLAG REGISTER (ADDRESS: 2Ch)

8.0 OSCILLATOR

The MCP2510 is designed to be operated with a crystal or ceramic resonator connected to the OSC1 and OSC2 pins. The MCP2510 oscillator design requires the use of a parallel cut crystal. Use of a series cut crystal may give a frequency out of the crystal manufacturers specifications. A typical oscillator circuit is shown in [Figure 8-1.](#page-48-0) The MCP2510 may also be driven by an external clock source connected to the OSC1 pin as shown in [Figure 8-2](#page-48-1) and [Figure 8-3.](#page-49-0)

8.1 Oscillator Startup Timer

The MCP2510 utilizes an oscillator startup timer (OST), which holds the MCP2510 in reset, to insure that the oscillator has stabilized before the internal state machine begins to operate. The OST maintains reset for the first 128 OSC1 clock cycles after power up, RESET, or wake up from sleep mode occurs. It should be noted that no SPI operations should be attempted until after the OST has expired.

8.2 CLKOUT Pin

The clock out pin is provided to the system designer for use as the main system clock or as a clock input for other devices in the system. The CLKOUT has an internal prescaler which can divide FOSC by 1, 2, 4 and 8. The CLKOUT function is enabled and the prescaler is selected via the CANCNTRL register (see [Register 9-](#page-51-0) [1](#page-51-0)). The CLKOUT pin will be active upon system reset and default to the slowest speed (divide by 8) so that it can be used as the MCU clock. When sleep mode is requested, the MCP2510 will drive sixteen additional clock cycles on the CLKOUT pin before entering sleep mode. The idle state of the CLKOUT pin in sleep mode is low. When the CLKOUT function is disabled (CAN-CNTRL.CLKEN = $'0'$ the CLKOUT pin is in a high impedance state.

The CLKOUT function is designed to guarantee that t_hCLKOUT and t_ICLKOUT timings are preserved when the CLKOUT pin function is enabled, disabled, or the prescaler value is changed.

FIGURE 8-1: CRYSTAL/CERAMIC RESONATOR OPERATION

FIGURE 8-3: EXTERNAL SERIES RESONANT CRYSTAL OSCILLATOR CIRCUIT

9.0 MODES OF OPERATION

The MCP2510 has five modes of operation. These modes are:

- 1. Configuration Mode.
- 2. Normal Mode.
- 3. Sleep Mode.
- 4. Listen-Only Mode.
- 5. Loopback Mode.

The operational mode is selected via the CANCTRL. REQOP bits (see [Register 9-1\)](#page-51-1). When changing modes, the mode will not actually change until all pending message transmissions are complete. Because of this, the user must verify that the device has actually changed into the requested mode before further operations are executed. Verification of the current operating mode is done by reading the CANSTAT. OPMODE bits (see [Register 9-2](#page-52-1)).

9.1 Configuration Mode

The MCP2510 must be initialized before activation. This is only possible if the device is in the configuration mode. Configuration mode is automatically selected after powerup or a reset, or can be entered from any other mode by setting the CANTRL.REQOP bits to '100'. When configuration mode is entered all error counters are cleared. Configuration mode is the only mode where the following registers are modifiable:

- CNF1, CNF2, CNF3
- TXRTSCTRL
- Acceptance Filter Registers
- Acceptance Mask Registers

Only when the CANSTAT.OPMODE bits read as '100' can the initialization be performed, allowing the configuration registers, acceptance mask registers, and the acceptance filter registers to be written. After the configuration is complete, the device can be activated by programming the CANCTRL.REQOP bits for normal operation mode (or any other mode).

9.2 Sleep Mode

The MCP2510 has an internal sleep mode that is used to minimize the current consumption of the device. The SPI interface remains active even when the MCP2510 is in sleep mode, allowing access to all registers.

To enter sleep mode, the mode request bits are set in the CANCTRL register. The CANSTAT.OPMODE bits indicate whether the device successfully entered sleep mode. These bits should be read after sending the sleep command to the MCP2510. The MCP2510 is active and has not yet entered sleep mode until these bits indicate that sleep mode has been entered. When in internal sleep mode, the wakeup interrupt is still active (if enabled). This is done so the MCU can also

be placed into a sleep mode and use the MCP2510 to wake it up upon detecting activity on the bus.

Note: Care must be exercised to not enter sleep mode while the MCP2510 is transmitting a message. If sleep mode is requested while transmitting, the transmission will stop without completing and errors will occur on the bus. Also, the message will remain pending and transmit upon wake up.

When in sleep mode, the MCP2510 stops its internal oscillator. The MCP2510 will wake-up when bus activity occurs or when the MCU sets, via the SPI interface, the CANINTF.WAKIF bit to 'generate' a wake up attempt (the CANINTF.WAKIF bit must also be set in order for the wakeup interrupt to occur). The TXCAN pin will remain in the recessive state while the MCP2510 is in sleep mode. Note that Sleep Mode will be entered immediately, even if a message is currently being transmitted, so it is necessary to insure that all TXREQ bits are clear before setting Sleep Mode.

9.2.1 WAKE-UP FUNCTIONS

The device will monitor the RXCAN pin for activity while it is in sleep mode. If the CANINTE.WAKIE bit is set, the device will wake up and generate an interrupt. Since the internal oscillator is shut down when sleep mode is entered, it will take some amount of time for the oscillator to start up and the device to enable itself to receive messages. The device will ignore the message that caused the wake-up from sleep mode as well as any messages that occur while the device is 'waking up.' The device will wake up in listen-only mode. The MCU must set normal mode before the MCP2510 will be able to communicate on the bus.

The device can be programmed to apply a low-pass filter function to the RXCAN input line while in internal sleep mode. This feature can be used to prevent the device from waking up due to short glitches on the CAN bus lines. The CNF3.WAKFIL bit enables or disables the filter.

9.3 Listen Only Mode

Listen-only mode provides a means for the MCP2510 to receive all messages including messages with errors. This mode can be used for bus monitor applications or for detecting the baud rate in 'hot plugging' situations. For auto-baud detection it is necessary that there are at least two other nodes, which are communicating with each other. The baud rate can be detected empirically by testing different values until valid messages are received. The listen-only mode is a silent mode, meaning no messages will be transmitted while in this state, including error flags or acknowledge signals. The filters and masks can be used to allow only particular messages to be loaded into the receive registers, or the filter masks can be set to all zeros to allow a message with any identifier to pass. The error

counters are reset and deactivated in this state. The listen-only mode is activated by setting the mode request bits in the CANCTRL register.

9.4 Loopback Mode

This mode will allow internal transmission of messages from the transmit buffers to the receive buffers without actually transmitting messages on the CAN bus. This mode can be used in system development and testing. In this mode the ACK bit is ignored and the device will allow incoming messages from itself just as if they were coming from another node. The loopback mode is a silent mode, meaning no messages will be transmitted while in this state, including error flags or acknowledge signals. The TXCAN pin will be in a reccessive state while the device is in this mode. The filters and masks can be used to allow only particular messages to be loaded into the receive registers. The masks can be set to all zeros to provide a mode that accepts all messages. The loopback mode is activated by setting the mode request bits in the CANCTRL register.

9.5 Normal Mode

This is the standard operating mode of the MCP2510. In this mode the device actively monitors all bus messages and generates acknowledge bits, error frames, etc. This is also the only mode in which the MCP2510 will transmit messages over the CAN bus.

REGISTER 9-1: CANCTRL - CAN CONTROL REGISTER (ADDRESS: XFh)

 $-n =$ Value at POR $'1' = B$ it is set $'0' = B$ it is cleared $x = B$ it is unknown

REGISTER 9-2: CANSTAT - CAN STATUS REGISTER (ADDRESS: XEh)

NOTES:

10.0 REGISTER MAP

The register map for the MCP2510 is shown in [Table 10-1.](#page-54-0) Address locations for each register are determined by using the column (higher order 4 bits) and row (lower order 4 bits) values. The registers have been arranged to optimize the sequential reading and

writing of data. Some specific control and status registers allow individual bit modification using the SPI Bit Modify command. The registers that allow this command are shown as shaded locations in [Table 10-1.](#page-54-0) A summary of the MCP2510 control registers is shown in [Table 10-2](#page-54-1).

Lower Address Bits	Higher Order Address Bits							
	x000 xxxx	$x001$ xxxx	$x010$ xxxx	x0011 xxxx	$x100$ xxxx	$x101$ xxxx	$x110$ $xxxx$	$x111$ $xxxx$
0000	RXF0SIDH	RXF3SIDH	RXMOSIDH	TXB0CTRL	TXB1CTRL	TXB2CTRL	RXB0CTRL	RXB1CTRL
0001	RXF0SIDL	RXF3SIDL	RXMOSIDL	TXB0SIDH	TXB1SIDH	TXB2SIDH	RXB0SIDH	RXB1SIDH
0010	RXF0EID8	RXF3EID8	RXM0EID8	TXB0SIDL	TXB1SIDL	TXB2SIDL	RXB0SIDL	RXB1SIDL
0011	RXF0EID0	RXF3EID0	RXM0EID0	TXB0EID8	TXB1EID8	TXB2EID8	RXB0EID8	RXB1EID8
0100	RXF1SIDH	RXF4SIDH	RXM1SIDH	TXB0EID0	TXB1EID0	TXB2EID0	RXB0EID0	RXB1EID0
0101	RXF1SIDL	RXF4SIDL	RXM1SIDL	TXB0DLC	TXB1DLC	TXB2DLC	RXB0DLC	RXB1DLC
0110	RXF1EID8	RXF4EID8	RXM1EID8	TXB0D0	TXB1D0	TXB2D0	RXB0D0	RXB1D0
0111	RXF1EID0	RXF4EID0	RXM1EID0	TXB0D1	TXB1D1	TXB2D1	RXB0D1	RXB1D1
1000	RXF2SIDH	RXF5SIDH	CNF ₃	TXB0D2	TXB1D2	TXB2D2	RXB0D2	RXB1D2
1001	RXF2SIDL	RXF5SIDL	CNF ₂	TXB0D3	TXB1D3	TXB2D3	RXB0D3	RXB1D3
1010	RXF2EID8	RXF5EID8	CNF ₁	TXB0D4	TXB1D4	TXB2D4	RXB0D4	RXB1D4
1011	RXF2EID0	RXF5EID0	CANINTE	TXB0D5	TXB1D5	TXB2D5	RXB0D5	RXB1D5
1100	BFPCTRL	TEC	CANINTF	TXB0D6	TXB1D6	TXB2D6	RXB0D6	RXB1D6
1101	TXRTSCTRL	REC	EFLG	TXB0D7	TXB1D7	TXB2D7	RXB0D7	RXB1D7
1110	CANSTAT	CANSTAT	CANSTAT	CANSTAT	CANSTAT	CANSTAT	CANSTAT	CANSTAT
1111	CANCTRL	CANCTRL	CANCTRL	CANCTRL	CANCTRL	CANCTRL	CANCTRL	CANCTRL

TABLE 10-1: CAN CONTROLLER REGISTER MAP

Note: Shaded register locations indicate that these allow the user to manipulate individual bits using the 'Bit Modify' Command.

TABLE 10-2: CONTROL REGISTER SUMMARY

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NOTES:

11.0 SPI INTERFACE

11.1 Overview

The MCP2510 is designed to interface directly with the Serial Peripheral Interface (SPI) port available on many microcontrollers and supports Mode 0,0 and Mode 1,1. Commands and data are sent to the device via the SI pin, with data being clocked in on the rising edge of SCK. Data is driven out by the MCP2510, on the SO line, on the falling edge of SCK. The \overline{CS} pin must be held low while any operation is performed. [Table 11-1](#page-57-0) shows the instruction bytes for all operations. Refer to [Figure 11-8](#page-59-0) and [Figure 11-9](#page-59-1) for detailed input and output timing diagrams for both Mode 0,0 and Mode 1,1 operation.

11.2 Read Instruction

The Read Instruction is started by lowering the $\overline{\text{CS}}$ pin. The read instruction is then sent to the MCP2510 followed by the 8-bit address (A7 through A0). After the read instruction and address are sent, the data stored in the register at the selected address will be shifted out on the SO pin. The internal address pointer is automatically incremented to the next address after each byte of data is shifted out. Therefore it is possible to read the next consecutive register address by continuing to provide clock pulses. Any number of consecutive register locations can be read sequentially using this method. The read operation is terminated by raising the \overline{CS} pin [\(Figure 11-2](#page-57-1)).

11.3 Write Instruction

The Write Instruction is started by lowering the $\overline{\text{CS}}$ pin. The write instruction is then sent to the MCP2510 followed by the address and at least one byte of data. It is possible to write to sequential registers by continuing to clock in data bytes, as long as \overline{CS} is held low. Data will actually be written to the register on the rising edge of the SCK line for the D0 bit. If the $\overline{\text{CS}}$ line is brought high before eight bits are loaded, the write will be aborted for that data byte, previous bytes in the command will have been written. Refer to the timing diagram in [Figure 11-3](#page-57-2) for more detailed illustration of the byte write sequence.

11.4 Request To Send (RTS) Instruction

The RTS command can be used to initiate message transmission for one or more of the transmit buffers.

The part is selected by lowering the \overline{CS} pin and the RTS command byte is then sent to the MCP2510. As shown in [Figure 11-4,](#page-58-0) the last 3 bits of this command indicate which transmit buffer(s) are enabled to send. This command will set the TxBnCTRL.TXREQ bit for the respective buffer(s). Any or all of the last three bits can be set in a single command. If the RTS command is sent with $n = 000$, the command will be ignored.

11.5 Read Status Instruction

The Read Status Instruction allows single instruction access to some of the often used status bits for message reception and transmission.

The part is selected by lowering the \overline{CS} pin and the read status command byte, shown in [Figure 11-6](#page-58-1), is sent to the MCP2510. After the command byte is sent, the MCP2510 will return eight bits of data that contain the status. If additional clocks are sent after the first eight bits are transmitted, the MCP2510 will continue to output the status bits as long as the $\overline{\text{CS}}$ pin is held low and clocks are provided on SCK. Each status bit returned in this command may also be read by using the standard read command with the appropriate register address.

11.6 Bit Modify Instruction

The Bit Modify Instruction provides a means for setting or clearing individual bits in specific status and control registers. This command is not available for all registers. See [Section 10.0](#page-54-2) (register map) to determine which registers allow the use of this command.

The part is selected by lowering the \overline{CS} pin and the Bit Modify command byte is then sent to the MCP2510. After the command byte is sent, the address for the register is sent followed by the mask byte and then the data byte. The mask byte determines which bits in the register will be allowed to change. A '1' in the mask byte will allow a bit in the register to change and a '0' will not. The data byte determines what value the modified bits in the register will be changed to. A '1' in the data byte will set the bit and a '0' will clear the bit, provided that the mask for that bit is set to a '1'. (see [Figure 11-1\)](#page-57-3)

11.7 Reset Instruction

The Reset Instruction can be used to re-initialize the internal registers of the MCP2510 and set configuration mode. This command provides the same functionality, via the SPI interface, as the RESET pin. The Reset instruction is a single byte instruction which requires selecting the device by pulling \overline{CS} low, sending the instruction byte, and then raising $\overline{\text{CS}}$. It is highly recommended that the reset command be sent (or the RESET pin be lowered) as part of the power-on initialization sequence. The MCP2510 will be held in reset for 128 FOSC cycles.

FIGURE 11-1: BIT MODIFY

TABLE 11-1: SPI INSTRUCTION SET

FIGURE 11-2: READ INSTRUCTION

FIGURE 11-3: BYTE WRITE INSTRUCTION

FIGURE 11-7: RESET INSTRUCTION

FIGURE 11-8: SPI INPUT TIMING

FIGURE 11-9: SPI OUTPUT TIMING

12.0 ELECTRICAL CHARACTERISTICS

12.1 Absolute Maximum Ratings†

† **Notice:** Stresses above those listed under "Maximum Ratings" may cause permanent damage to the device. This is a stress rating only and functional operation of the device at those or any other conditions above those indicated in the operational listings of this specification is not implied. Exposure to maximum rating conditions for extended periods may affect device reliability.

TABLE 12-1: DC CHARACTERISTICS

Note: This parameter is periodically sampled and not 100% tested.

TABLE 12-2: OSCILLATOR TIMING CHARACTERISTICS

Note: This parameter is periodically sampled and not 100% tested.

TABLE 12-3: CAN INTERFACE AC CHARACTERISTICS

TABLE 12-4: CLKOUT PIN AC/DC CHARACTERISTICS

Note: CLKOUT prescaler set to divide by one.

TABLE 12-5: SPI INTERFACE AC CHARACTERISTICS

Note: This parameter is not 100% tested.

13.0 PACKAGING INFORMATION

13.1 Package Marking Information

18-Lead Plastic Dual In-Line (P) – 300 mil Body [PDIP]

Note: For the most current package drawings, please see the Microchip Packaging Specification located at http://www.microchip.com/packaging

Notes:

- 1. Pin 1 visual index feature may vary, but must be located within the hatched area.
- 2. § Significant Characteristic.
- 3. Dimensions D and E1 do not include mold flash or protrusions. Mold flash or protrusions shall not exceed .010" per side.
- 4. Dimensioning and tolerancing per ASME Y14.5M.

BSC: Basic Dimension. Theoretically exact value shown without tolerances.

Microchip Technology Drawing C04-007B

18-Lead Plastic Small Outline (SO) – Wide, 7.50 mm Body [SOIC]

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Notes:

- 1. Pin 1 visual index feature may vary, but must be located within the hatched area.
- 2. § Significant Characteristic.
- 3. Dimensions D and E1 do not include mold flash or protrusions. Mold flash or protrusions shall not exceed 0.15 mm per side.
- 4. Dimensioning and tolerancing per ASME Y14.5M.
	- BSC: Basic Dimension. Theoretically exact value shown without tolerances.

REF: Reference Dimension, usually without tolerance, for information purposes only.

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20-Lead Plastic Thin Shrink Small Outline (ST) – 4.4 mm Body [TSSOP]

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Notes:

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2. Dimensions D and E1 do not include mold flash or protrusions. Mold flash or protrusions shall not exceed 0.15 mm per side.

- 3. Dimensioning and tolerancing per ASME Y14.5M.
	- BSC: Basic Dimension. Theoretically exact value shown without tolerances.

REF: Reference Dimension, usually without tolerance, for information purposes only.

Microchip Technology Drawing C04-088B

APPENDIX A: REVISION HISTORY

Revision F (January 2007)

This revision includes updates to the packaging diagrams.

NOTES:

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